



OH 4(1) 2021 AA. IT.3-3.

$n_i = \# \text{ nodes in heap } D_i \text{ after } i \text{ operations}$   
→ natural log

$$\phi(D_i) = K n_i \ln(n_i)$$

$K = \text{large enough const}$   
based on heap ops ins ext-min

Trick  $n \cdot \ln\left(\frac{n}{n-1}\right) = \Theta(1)$

$$\Leftrightarrow n \cdot \ln\left(\frac{n}{n-1}\right) \leq \text{const} + C$$

$$\text{op} = \Theta(\log n) \Rightarrow K \text{ const}$$

$$\Leftrightarrow \lim_{n \rightarrow \infty} n \ln\left(\frac{n}{n-1}\right) = \text{finite const}$$

proof:  $\lim_{n \rightarrow \infty} n \ln \frac{n}{n-1} = \lim_{n \rightarrow \infty} \ln \left[ \left(1 + \frac{1}{n-1}\right)^n \right] =$

$$= \lim_{n \rightarrow \infty} \ln \left[ \underbrace{\left(1 + \frac{1}{n-1}\right)^{n-1}}_{\lim_n(1) = e} \right]^{n-1} = \lim_{n \rightarrow \infty} \ln(e) =$$
$$= \lim_{n \rightarrow \infty} 1^{\frac{n-1}{n-1}} = 1$$

INSERT

$$n_{i-1} = h_i - 1$$

$$\hat{c}_i = c_i + \phi(\delta_i) - \phi(\delta_{i-1})$$

$$\leq k \ln(n_i) \text{ true cost upper bound}$$

$$\leq k \ln(n_i) + k n_i \ln(n_i) - k n_{i-1} \ln(n_{i-1})$$

$$= k \ln(n_i) + k n_i \ln(n_i) - k^{h_i-1} \ln(n_i-1)$$

exclude

$$\leq k \ln(n_i) + k n_i \ln(n_i) - k^{h_i} \ln(n_i-1)$$

$$= k \ln(n_i) + k n_i \ln(n_i) - k^{h_i} \ln(n_i-1) + k \ln(n_i-1)$$

$$\leq 2k \ln(n_i) + k n_i [\ln(n_i) - \ln(n_i-1)]$$

$$= 2k \ln(n_i) + k n_i \ln \frac{n_i}{n_i-1} \text{ trick } \leq \text{const } C$$

$$\leq 2k \ln(n_i) + k \cdot C \quad k, C \text{ constants}$$

$$= O(\log(n_i))$$

$$\boxed{\text{Extract Min}} \quad \hat{G} = c_i + \phi(D) - \phi(D_{i-1})$$

$$n_i = n_{i-1} - 1$$

$$\leq k \ln(n_i)$$

$$\leq k \ln(n_{i-1}) + k n_i \ln(n_i) - k n_{i-1} \ln(n_{i-1})$$

$$= k \ln(n_{i-1}) + k n_i \ln(n_{i-1} - 1) - k n_{i-1} \ln(n_{i-1})$$

$$= k \ln(n_{i-1}) + \cancel{k n_{i-1} \ln(n_{i-1} - 1)} \rightarrow -k \ln(n_{i-1} - 1) - \cancel{k n_{i-1} \ln(n_{i-1})}$$

$$= k \left[ \ln(n_{i-1}) - \ln(n_{i-1} - 1) \right] + \cancel{k n_{i-1} \left[ \ln(n_{i-1} - 1) - \ln(n_{i-1}) \right]}$$

$$= \dots - \ln\left(\frac{n_{i-1}}{n_{i-1}-1}\right)$$

trick

$$= \dots$$

$$= O(1)$$

Easier  $\phi(O_i)$  :

(hard to come up with)  $x = \text{node/value in heap}$   
 $\text{depth}_i(x) = \text{Depth/level of } x$   
after  $i$ th operation.

$$\phi(O_i) = \sum_{x \in \text{heap}} k(\text{depth}(x) + 1)$$

→ because each op only affects one value ( $x$ )