

# Optimize Your Programs, Elegantly

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# Agenda

- 🕒 Introduction.
- 🕒 Dijkstra's Shortest Paths.
- 🕒 Topological Ordering.

# You can make your programs more elegant by

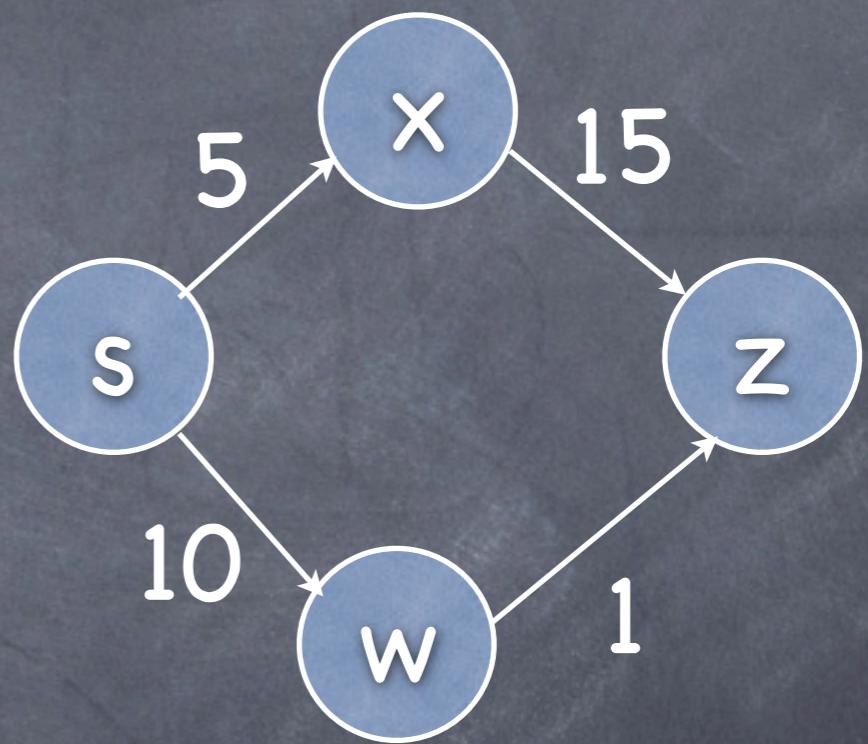
- ⦿ Having “components that influence other components”.
- ⦿ BTW, these are called “aspects”.
- ⦿ AspectJ is a general purpose programming languages for developing aspects.
- ⦿ You can also simulate some what AspectJ aspects do using inheritance:
  - ⦿ A subclass is an influenced version of its superclass. (Java)
  - ⦿ Mixins, Traits produce influenced versions of classes. (Scala)

# Figuring Out Aspects

- ⦿ We recommend that you figure out aspects right from the requirements rather than refactoring your program into aspects.
- ⦿ Efficiency is one such requirement that can be implemented as an aspect.

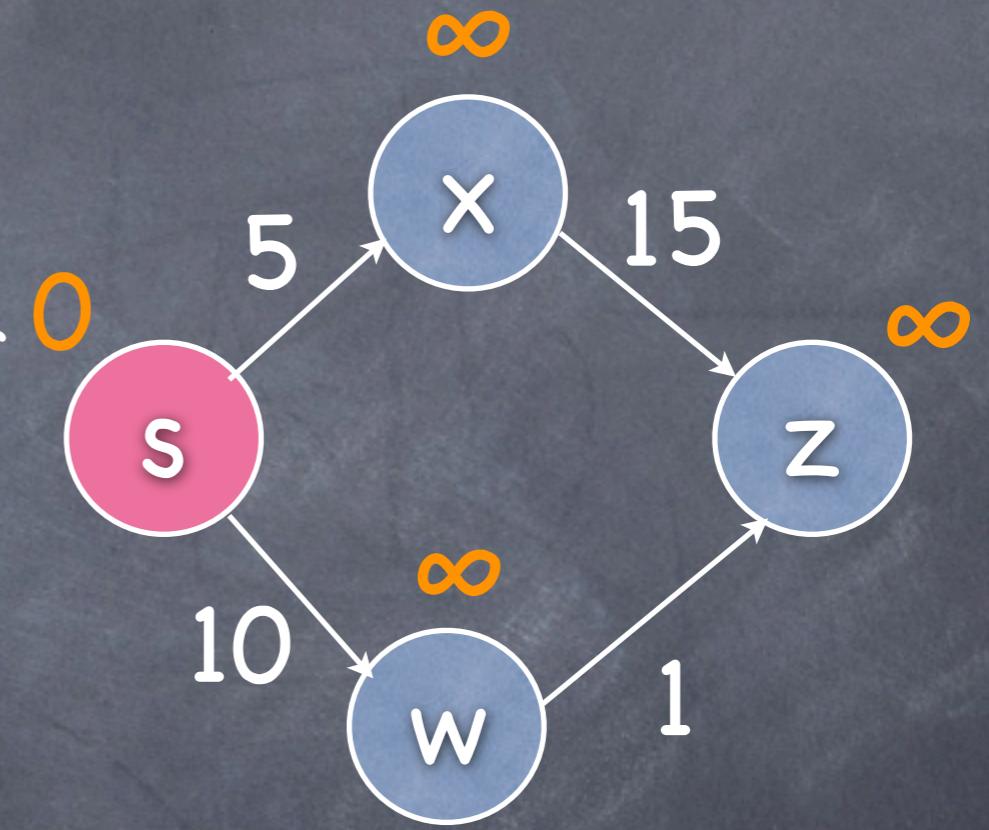
# Dijkstra's Shortest Paths

- Given a weighted Graph (g) and a source node (s), find the shortest path from s to all nodes in g.



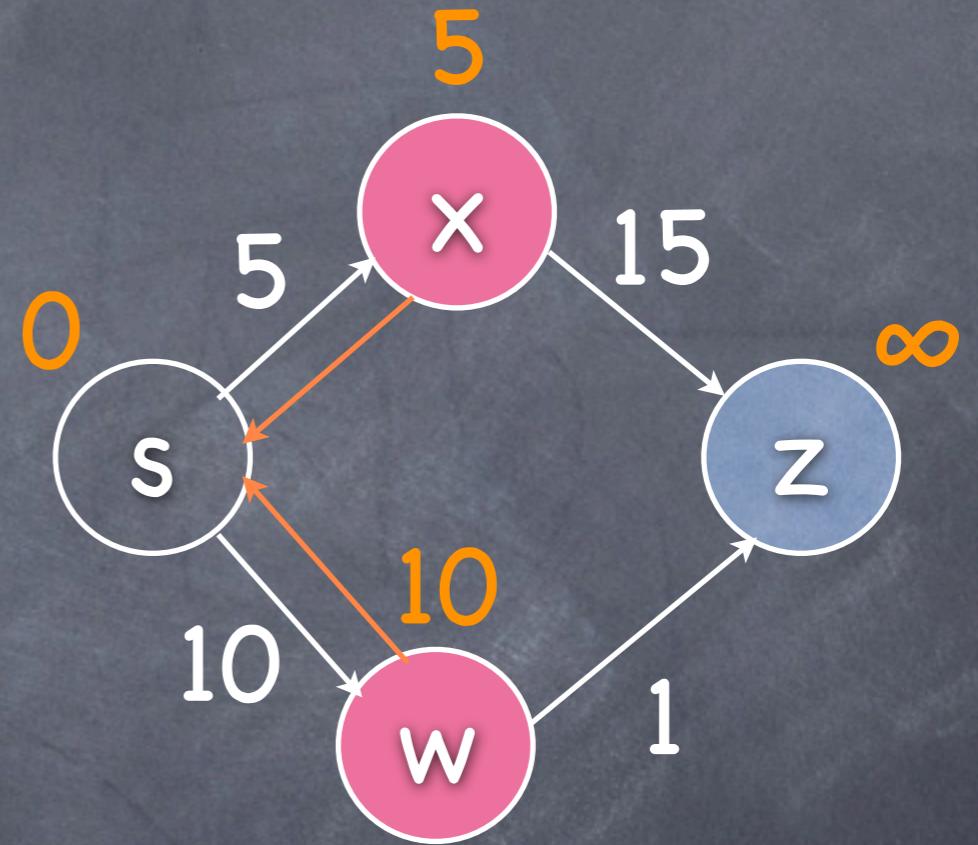
# Dijkstra's Shortest Paths

- To each node, maintain its best known distance from source and its predecessor on the shortest path from source.
- Maintain a list of active nodes. Initially = {S}.
- Maintain a list of final nodes. Initially = {}



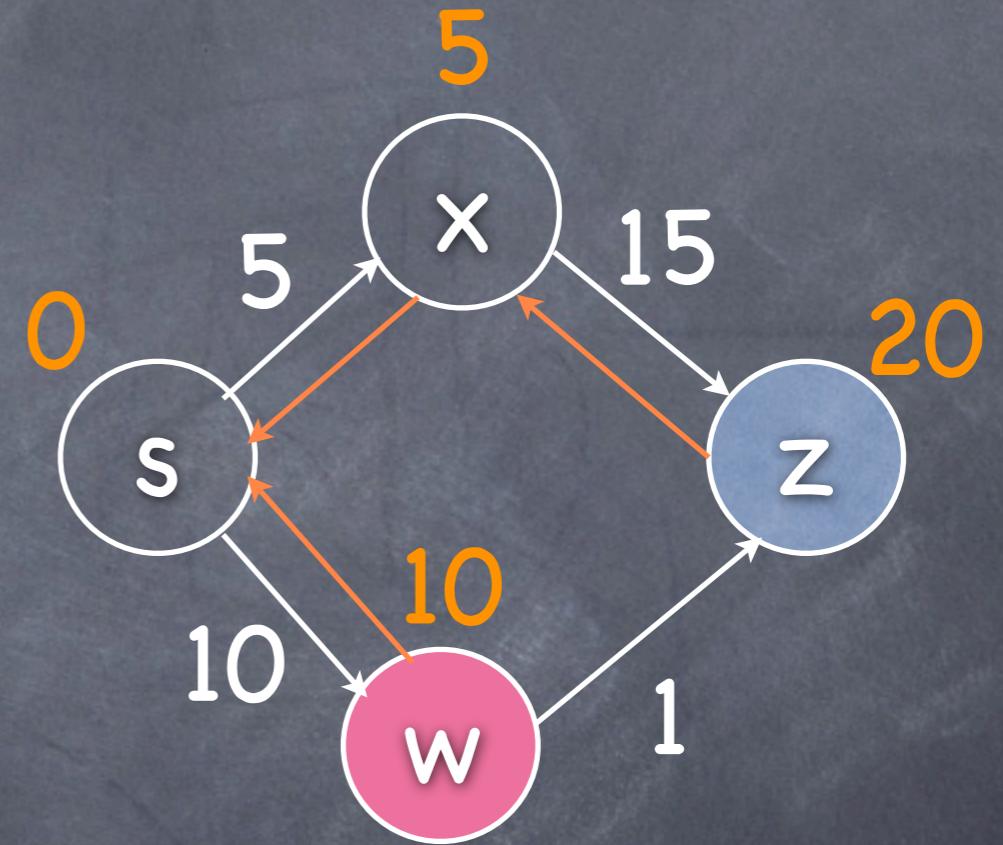
# Dijkstra's Shortest Paths

- Select the closest to source active node ( $v$ ).
- Update the best known distance as well as predecessors of  $v$ 's successors that are not marked as final.
- Mark  $v$  as final.



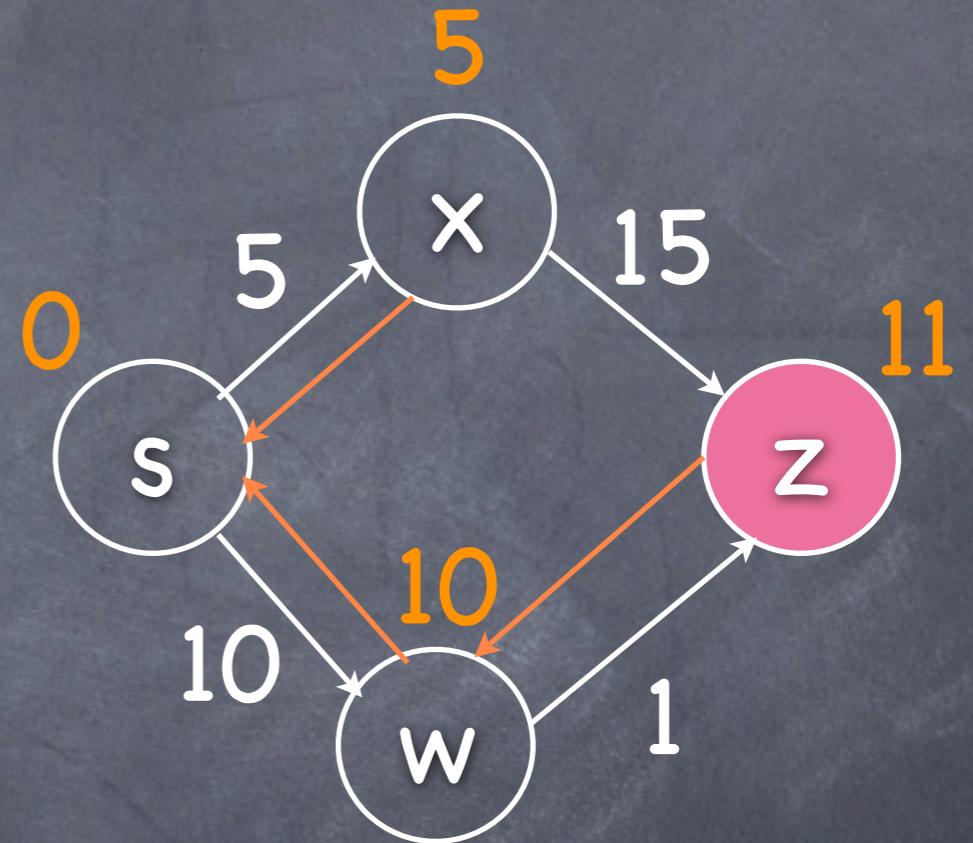
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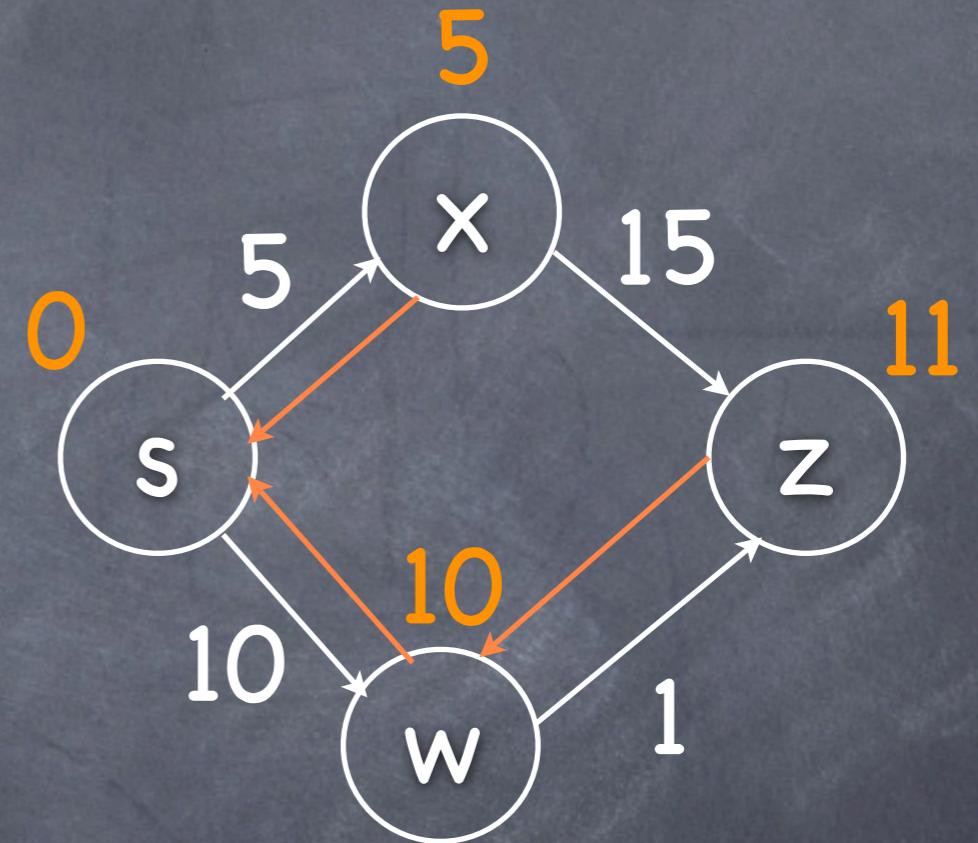
# Dijkstra's Shortest Paths

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# Dijkstra's Shortest Paths

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- Update the best known distance as well as predecessors of  $v$ 's successors that are not marked as final.
- Mark  $v$  as final.



# Straight Forward Implementation (1)

To node we add:

```
private int distance = Integer.MAX_VALUE;
private Node predecessor = null;
private boolean finalized = false;

public void init(boolean isSource){
    finalized = false;
    predecessor = null;
    distance = isSource? 0 :Integer.MAX_VALUE;
}

public boolean isFinal(){ return finalized; }

public void markAsFinal(){ finalized = true; }

public boolean isReachable(){
    return distance != Integer.MAX_VALUE; }

public boolean isActive(){
    return !isFinal() && isReachable();}
```

```
public Maybe<Path> getShortestPath(){...}

public void updateDistances(){
    for (Edge e : getOutgoingEdges()) {
        Node v = (Node) e.getTarget();
        if(!v.isFinal()){
            int alt = distance + e.getLength();
            if(alt < v.distance){
                v.distance = alt;
                v.predecessor = this;
            }
        }
    }
}

@Override
public int compareTo(Node o) {
    return distance - o.distance;
}
```

# Straight Forward Implementation (2)

- Note that:
  - Instead of maintaining the set of final nodes, we maintain a finalized bit for each node. Otherwise, we need to pass too much information (the list of all final nodes) to `Node.updateDistances()`.
  - `distance`, `finalized` and `predecessor` are only meaningful during an invocation of `dijkstra(..)`.

# Straight Forward Implementation (3)

To graph we add:

```
public Map<Node, Maybe<Path>> dijkstra(Node source){  
    // Initialize Dijkstra node extension  
    for (Node n : getNodes()) {  
        n.init(n == source);  
    }  
    // Compute shortest paths  
    for(Node u = getClosestActiveNode();  
        u != null;  
        u = getClosestActiveNode()){  
        u.markAsFinal();  
        u.updateDistances();  
    }  
    // retrieve shortest paths  
    Map<Node, Maybe<Path>> result = new ...;  
    for (Node node : getNodes()) {  
        result.put(node, node.getShortestPath());  
    }  
    return result;  
}
```

```
public Node getClosestActiveNode(){  
    Collection<Node> activeNodes = getActiveNodes();  
    if(activeNodes.size() == 0){  
        return null;  
    }else{  
        return Collections.min(activeNodes);  
    }  
  
private Collection<Node> getActiveNodes(){  
    List<Node> active = new ...;  
    for (Node n : getNodes()) {  
        if(n.isActive()){  
            active.add(n);  
        }  
    }  
    return active;  
}
```

```
32:36:1152:public dijkstraaj.Graph$Node dijkstraaj.Graph.getClosestNonFinalNode()[0]
30:36:1080:public java.util.Collection dijkstraaj.Graph.getActiveNodes()[0]
```

# Optimizations

- ⦿ Analyzing the execution traces of our implementation we discover that the two methods `getClosestActiveNode()` and `getActiveNodes()` are invoked several times with the same graph object.
- ⦿ We maintain a priority queue of nodes ordered by the node's shortest known distance from source and use it to speed up `getClosestActiveNode()`.

# Optimization Aspect

```
public aspect ActiveNodes
percflow(execution(public Map<Node, Maybe<Path>> dijkstra(Node))){
    //active nodes in the graph ordered by their distance from source
    private PriorityQueue<Node> activeNodes = new ...;

    after(Node n): call(public void markAsFinal()) && target(n){
        activeNodes.remove(n);
    }

    before(Node n): set(private int Node.distance) && target(n){
        if(n.isReachable()){
            activeNodes.remove(n);
        }
    }

    after(Node n): set(private int Node.distance) && target(n){
        if(n.isReachable()){
            activeNodes.add(n);
        }
    }

    Node around(): call(public Node getClosestActiveNode()){
        return activeNodes.peek();
    }
}
```

# Why keeping the optimization separate?

- ⦿ No need to change method signatures (what's the problem with changing method signatures?).
- ⦿ Fewer things to keep in your head while:
  - ⦿ Developing the algorithm.
  - ⦿ Understanding/Maintaining the algorithm.
- ⦿ Reduced scattering and tangling of optimization related code with the base program.
- ⦿ Safety: Algorithm + Optimization Aspect ≈ Algorithm

# Changing Method Signatures

```
public Map<Node, Maybe<Path>> dijkstra(Node source){  
=> private PriorityQueue<Node> activeNodes = new ...;  
    // Initialize Dijkstra node extension  
    for (Node n : getNodes()) {  
        n.init(n == source);  
    }  
    // Compute shortest paths  
=> for(Node u = activeNodes.peek();  
        u != null;  
=>     u = activeNodes.peek()){  
        u.markAsFinal();  
=>        if(v.isReachable()){  
            activeNodes.remove(u);  
=>        }  
*=>        u.updateDistances(activeNodes);  
    }  
    // retrieve shortest paths  
    Map<Node, Maybe<Path>> result = new ...;  
    for (Node node : getNodes()) {  
        result.put(node, node.getShortestPath());  
    }  
    return result;  
}
```

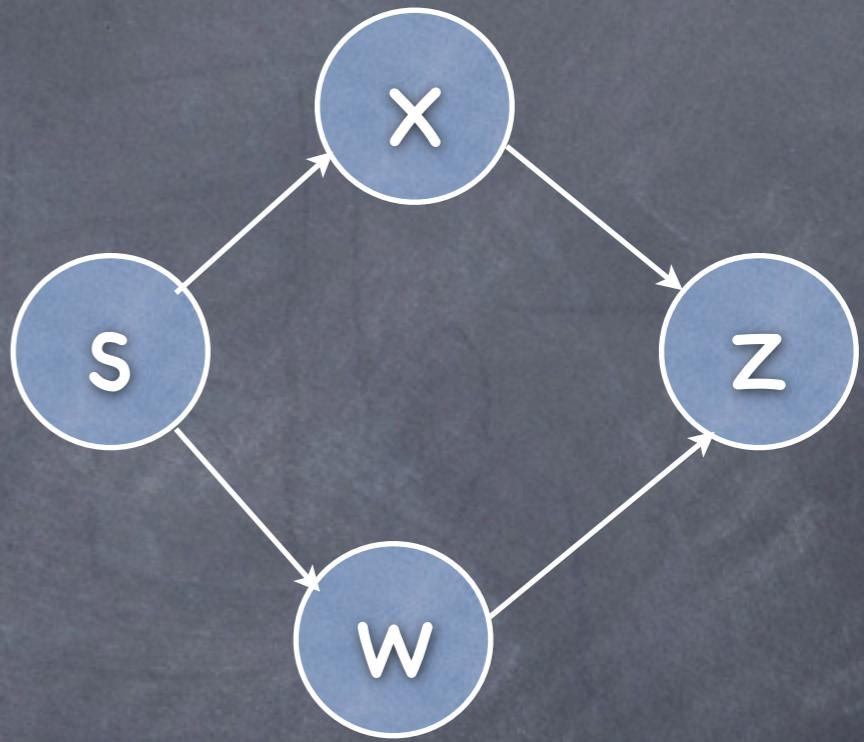
```
=> public void updateDistances  
(PriorityQueue<Node> activeNodes){  
    for (Edge e : getOutgoingEdges()) {  
        Node v = (Node) e.getTarget();  
        if(!v.isFinal()){  
            int alt = distance + e.getLength()  
            if(alt < v.distance){  
                if(v.isReachable()){  
                    activeNodes.remove(v);  
                }  
                v.distance = alt;  
                if(v.isReachable()){  
                    activeNodes.add(v);  
                }  
                v.predecessor = this;  
            }  
        }  
    }  
}  
=> public void init(boolean isSource,  
PriorityQueue<Node> activeNodes){  
    finalized = false;  
    predecessor = null;  
=> if(isReachable()){  
    activeNodes.remove(this);  
=> }  
    distance = isSource? 0 :Integer.MAX_VALUE;  
=> if(isReachable()){  
    activeNodes.add(this);  
=> }  
}
```

# Safety

- ⦿ Algorithm + Optimization Aspect  $\approx$  Algorithm
- ⦿ What if you forgot one of the lines related to optimization?
- ⦿ Unit tests can catch the problem, but you'll get a single bit of information.
- ⦿ We have a tool that monitors the execution of optimization aspects and notifies you right after the optimization aspects first derails the algorithm rather than when the algorithm finishes.

# Topological Ordering

- Given a graph (g), find an ordering of its nodes such that node (a) comes before node (b) if there is a path from a to b.
- $\langle S, X, Z, W \rangle$  ?
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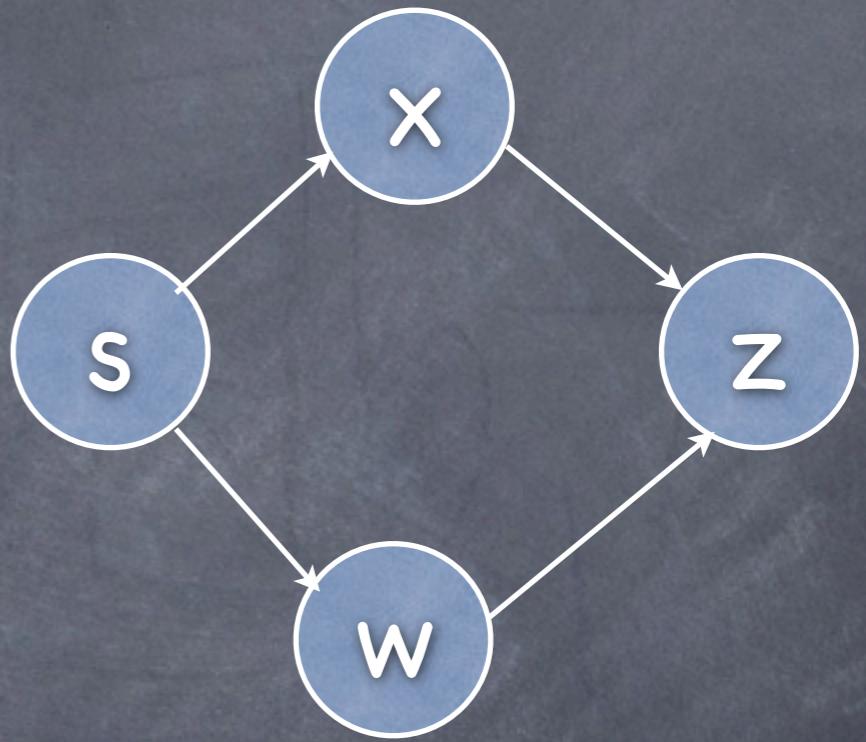


# Topological Ordering

```
public Maybe<List<Node>> topord(){
    List<Node> orderedNodes = new ...;
    for(Maybe<Node> mbNode = getSourceNode();
        mbNode.isSome();
        mbNode = getSourceNode()){
        orderedNodes.add(mbNode.get());
        remove(mbNode.get());
    }
    if(getNumNodes() > 0 ){
        return new None<List<Node>>();
    }
    return new Some<List<Node>>(orderedNodes);
}
```

```
public Maybe<Node> getSourceNode(){
    for (Node n : getNodes()) {
        if(n.getPredCount() == 0){
            return new Some<Node>(n);
        }
    }
    return new None<Node>();
}
```

```
public int getPredCount(){
    int predCount = 0;
    for (Node n : getEnclosingGraph().getNodes()) {
        if(n.hasSuccessor(this)) predCount++;
    }
    return predCount;
}
```

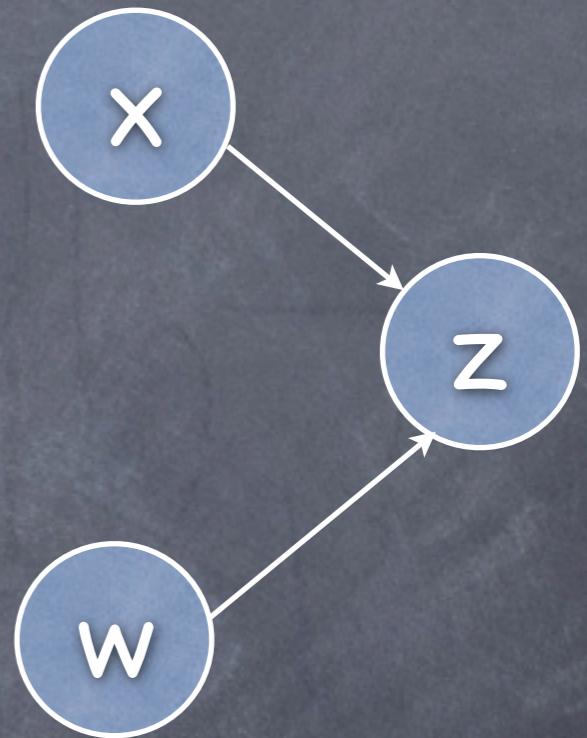


# Topological Ordering

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public Maybe<List<Node>> topord(){
    List<Node> orderedNodes =
        new ArrayList<Node>();
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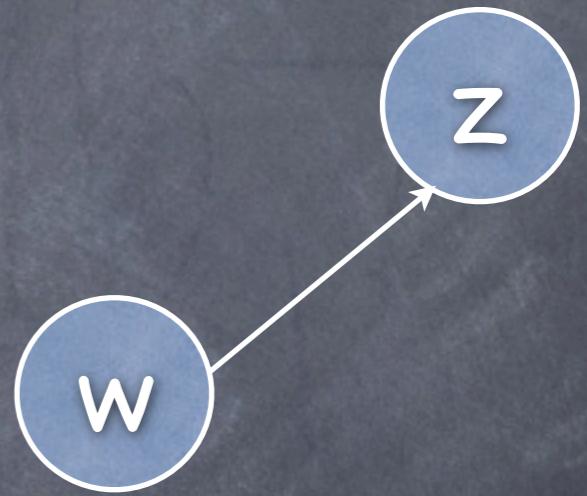
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public int getPredCount(){
    int predCount = 0;
    for (Node n : getEnclosingGraph().getNodes()) {
        if(n.hasSuccessor(this)) predCount++;
    }
    return predCount;
}
```

- ⦿ The graph is gone!
  - ⦿ Mark nodes as removed rather than actually removing them.
  - ⦿ Graph accessors need to be modified to take those marks into account during (in the control flow of) the execution of topord()
  - ⦿ Sounds like an aspect. Described as a modification to an existing component (the graph).

```
70:33:2310:public lib.Maybe topsort.Graph.getSourceNode()[0]  
29:77:2233:public int topsort.Graph$Node.getPredCount()[0]
```

# Not Optimal

- ⦿ getSourceNode() and getPredCount() are invoked several times with the same arguments.
- ⦿ Maintain the list of current source nodes. Upon removing a node
  - ⦿ remove it from the list.
  - ⦿ potentially add its successors to the list.
- ⦿ Maintain the count of predecessors. Upon removing a node, decrement the predecessors count of its successors.

# getSourceNode()

```
public aspect SrcNode {  
    private List<Node> Graph.sourceNodes = null;  
    //Memoization + Eager initialization  
    Maybe<Node> around(Graph g): execution(public Maybe<Node> getSourceNode()) && target(g){  
        if(g.sourceNodes == null){  
            g.sourceNodes = new ArrayList<Node>();  
            for (Node n : g.getNodes()) {  
                if(n.getPredCount() == 0){  
                    g.sourceNodes.add(n);  
                }  
            }  
        }  
        if(g.sourceNodes.size() == 0){  
            return new None<Node>();  
        }else{  
            return new Some<Node>(g.sourceNodes.get(0));  
        }  
    }  
    //Maintenance  
    void around(Graph g, Node n): execution(public void remove(Node)) && target(g) && args(n){  
        List<Node> successors0fToBeRemovedNode = new ArrayList<Node>(n.getSuccessors());  
        proceed(g,n);  
        for (Node succ : successors0fToBeRemovedNode) {  
            if(succ.getPredCount() == 0){  
                g.sourceNodes.add(succ);  
            }  
        }  
        g.sourceNodes.remove(n);  
    }  
}
```

# getPredCount()

```
public aspect Preds {  
  
    private int Node.predCount = -1;  
  
    //Memoization + Eager initialization  
    int around(Node node): execution(public int getPredCount(..)) && target(node){  
        if(node.predCount == -1){  
            for (Node n : node.getEnclosingGraph().getNodes()) {  
                n.predCount = 0;  
            }  
            for (Node n : node.getEnclosingGraph().getNodes()) {  
                for (Node succ : n.getSuccessors()) {  
                    succ.predCount++;  
                }  
            }  
        }  
        return node.predCount;  
    }  
  
    //Maintenance  
    before(Node n): execution(private void removeAllSuccessors(..)) && target(n){  
        for (Node succ : n.getSuccessors()) {  
            succ.predCount--;  
        }  
    }  
}
```

# To Take Home

- ⦿ Factor your program into components and aspects.
- ⦿ Don't refactor your program into aspects, it's harder that way.
- ⦿ Figure out aspects from the requirements.
- ⦿ Optimizations should be factored out into aspects.