

CS3000: Algorithms & Data

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Lecture 11:

- Midterm review

Feb 13, 2019

Course Contents

- Stable Matching Problem
 - Choosing definitions of Stability
 - Found counter examples of stability
 - Proved that some outcome is or is not stable
 - Gale Shapley Algorithm
 - Proved correctness via Induction
 - Computed the run time

- Asymptotics

- Asymptotics

- Way that we evaluate algorithms
- Only care about behavior for large n (limit as $n \rightarrow \infty$)
- Theta, Omega, Big-Oh, little omega, little oh
- All exponential grow faster than all polynomials
- All polynomials grow faster than logarithms to any power

- Induction

- Prove mathematical facts
- Used this to prove correctness of algorithms

- Divide and Conquer Algorithms
 - Merge Sort
 - Binary Search in a sorted list
 - Max Subarray Sum
 - Karatsuba

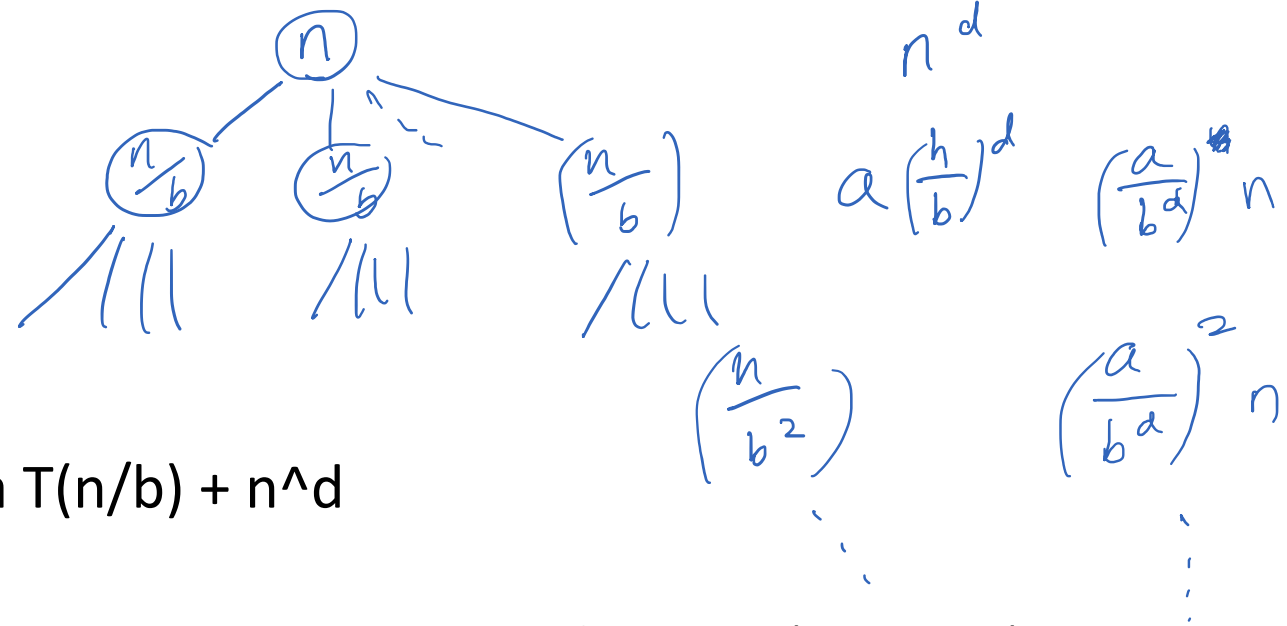
- Sorting

- Insertion Sort $\Theta(n^2)$
- Can sort in $n \log(n)$ time. Via MergeSort
- Sorting makes searching easier (via binary search)

Karatsuba

$$T(n) = 3 \cdot T\left(\frac{n}{2}\right) + cn$$

$$\left(\frac{3}{2^2}\right) = \frac{3}{2} > 1 \quad O(n^{\log_2 3})$$



• Master Theorem

• Solution to the recurrence $T(n) = a T(n/b) + n^d$

• Three cases:

• $(a/b^d) < 1$. $T(n) = \Theta(n^d)$ Work is dominated at first layer (initial call)

• $(a/b^d) = 1$ $T(n) = \Theta(\log(n) * n^d)$ Work at each layer is equal

→ • $(a/b^d) > 1$ $T(n) = \Theta(a^{\log_b n} = n^{\log_b a})$. Work is dominated at last layer (base cases)

• Where In the recursion tree was the work done?

• Number of layers $\log_b(n)$

- Create Algorithms
- Proved Correctness
- Analyzed Runtimes
- Found examples that cause algorithms to fail
- Write Pseudocode

Which is asymptotically smaller

~~2ⁿ~~ 2ⁿ or 3ⁿ

Can we say $2^n = o(3^n)$? YES

$$\lim_{n \rightarrow \infty} \frac{2^n}{3^n} = \lim_{n \rightarrow \infty} \left(\frac{2}{3}\right)^n = 0 \quad \checkmark$$

$$Is \quad 2^{n/2} = o(2^n) ?$$

$$\lim_{n \rightarrow \infty} \frac{2^{n/2}}{2^n} = \lim_{n \rightarrow \infty} 2^{-n/2} = 0$$

$$2^n \quad \text{vs} \quad 3^{n/2}$$
$$e^{\underbrace{n \lg 2}} \quad e^{\underbrace{\frac{n}{2} \lg 3}}$$

$$a^b = e^{b \ln a}$$