## CS3000: Algorithms \& Data Paul Hand

## Day 1:

- Course Overview
- Warmup Exercise (Induction, Asymptotics, Fun)

Jan 7, 2019

## Instructor

- Name: Paul Hand
- Call me Paul
- NEU since Fall 2018
- Office: 523 Lake
- Office Hours: Mon 1:15-2:45

- Research:
- Machine Learning, Artificial Intelligence, Computer Vision
- Algorithms are at the core of all of these!


## Discussion:

What impressive things can computers now do (but couldn't when you were born)?

What do you think computers will be able to do in 10 years that they can't today?

## Discussion:

What would you say is driving computational advances?

## Algorithms

- What is an algorithm?

An explicit, precise, unambiguous, mechanicallyexecutable sequence of elementary instructions for solving a computational problem.
-Jeff Erickson

- Essentially all computer programs (and more) are algorithms for some computational problem.


## Algorithms

- What is Algorithms?

The study of how to solve computational problems.

- Abstract and formalize computational problems
- Identify broadly useful algorithm design principles for solving computational problems
- Rigorously analyze properties of algorithms
- correctness, running time, space usage



## Moore's Law

Moore's Law - The number of transistors on integrated circuit chips (1971-2016)

## This advancement is important as other aspects of technological progress - such as processing speed or the price of electronic products - are

 strongly linked to Moore's law.

| News | Opinion | Sport | Culture | Lifestyle | More |
| :--- | :--- | :--- | :--- | :--- | :--- |

US World Environment Soccer US Politics Business Tech Science

## Artificial intelligence (AI)

## AlphaGo seals 4-1 victory over Go grandmaster Lee Sedol

## DeepMind's artificial intelligence astonishes fans to defeat human

 opponent and offers evidence computer software has mastered a major challenge
## Steven Borowiec



## A computational problem

- Each of you:

Determine how many other people in this room have the same first name as you.

## There are multiple ways to solve problems

- Sometimes your first instinct approach is reasonable.
- Sometimes your first instinct is not.


## Example problems we will look at

 $\bullet$
-



Finding the shortest path
helix HO

- 00000000

RYDSRTTIFSP. .EGRLYQVEYAMEAIGNA. GSATGILS RYDSRTTIFSPLREGRLYQVEYAMEAISHA.GTCLGILS RYDSRTTIFSP. .EGRLYQVEYAQEAISNA.GTAIGILS RYDSRTTIFSP. .EGRL YQVEYAMEAISHA.GTCLGILA RYDSRTTIFSP. .EGRL YQVEYAMEAIGHA.GTCLGILA RYDSRTTIFSP. .EGRLYQVEYAMEAIGNA. GSALGVLA RYDSRTTTESP. EGRLYQVEYALEAINNA. SITIGLIT SYDSRTTIFSP. . EGRLYQVEYALEAINHA.GVALGIVA

$$
(F, Y \text { or } W)_{15} S_{16} P_{17}
$$

Sequence Alignment

## Other good problems (we wont see)



Planted Clique

## Algorithms

- What is CS3000: Algorithms \& Data?

The study of how to solve computational problems. How to rigorously prove properties of algorithms.

- Proofs are about understanding and communication, not about formality or certainty
- Different emphasis from courses on logic
- We'll talk a lot about proof techniques and what makes a correct and convincing proof


## Algorithms

- That sounds hard. Why would I want to do that?
- Build Intuition:
- How/why do algorithms really work?
- How to attack new problems?
- Which design techniques work well?
- How to compare different solutions?
- How to know if a solution is the best possible?


## Algorithms

- That sounds hard. Why would I want to do that?
- Improve Communication:
- How to explain solutions?
- How to convince someone that a solution is correct?
- How to convince someone that a solution is best?


## Algorithms

- That sounds hard. Why would I want to do that?
- Get Rich:
- Many of the world's most successful companies (eg. Google) began with algorithms.
- Many job interviews have algorithm questions
- Understand the natural world:
- Brains, cells, networks, etc. often viewed as algorithms.
- Fun:
- Yes, seriously, fun.


## Course Structure



- $\mathrm{HW}=45 \%$
- Exams = 55\%
- Midterm I = 15\%
- Midterm II = 15\%
- Final = 25\%


## Course Structure



## Textbook:

Algorithm Design by Kleinberg and Tardos

More resources on the course website

## The TA Team

- TBD
- Office Hours: TBD
- Location: TBD
- TBD
- Office Hours: TBD
- Location: TBD
- TBD
- Office Hours: TBD
- Location: TBD


## Homework

- Weekly HW Assignments (45\% of grade)
- Due Wednesdays by 2:50pm
- HW1 out on Wednesday! Due Wed 1/16
- No extensions, no late work
- Lowest HW score will be dropped from your grade
- A mix of mathematical and algorithmic questions


## Homework Policies

- Homework must be typeset in LaTeX!
- Many resources available
- Many good editors available (TexShop, TexStudio)
- I will provide HW source

The Not So Short
Introduction to LATEX $2 \varepsilon$

Or $A A T_{E} X 2 \varepsilon$ in 157 minutes

## Homework Policies

- Homework will be submitted on Gradescope!
- More details on Wednesday
ill gradescope


## Homework Policies

- You are encouraged to work with your classmates on the homework problems.
- You may not use the internet
- You may not use students/people outside of the class
- Collaboration Policy:
- You must write all solutions by yourself
- You may not share any written solutions
- You must state all of your collaborators
- We reserve the right to ask you to explain any solution


## Discussion Forum

- We will use Piazza for discussions
- Ask questions and help your classmates
- Please use private messages sparingly
- More details on Wednesday!



## Course Website

## http://www.ccs.neu.edu/home/hand/teaching/cs3000-spring-2018/

## ccs.neu.edu



## CS3000: Algorithms \& Data

## Syllabus $\quad$ Schedule

This schedule will be updated frequently -check back often!

| \# | Date | Topic | Reading | HW |
| :--- | :--- | :--- | :--- | :--- |
| 1 | M 1/7 | Course Overview, Induction <br> Slides: | --- |  |
| 2 | W 1/9 | Stable Matching: Gale-Shapley Algorithm, Proof <br> by Contradiction | KT 1.1,1.2,2.3 | HW1 Out (pdf, <br> tex) |
| Slides: |  |  |  |  |

## What About the Other Sections?

- Prof. Schnyder teaches another section
- No formal relationship with my section
- Will cover very similar topics and share some materials
- Will be out of sync
- You should not go to OH for Prof. Schnyder's TAs


## Illustration: Let's count how many students are in this class

## Simple Counting

SimCount:
Find first student
First student says 1
Until we're out of students:
Go to next student
Next student says (what last student said + 1)

- Is this correct?
- How long does this take with n students?


## Recursive Counting - Divide and Conquer

RecursiveCount:
If you are the only person in group: return 1
Else:
Split your group into two subgroups of similar size (one includes you)
Appoint a leader of the other subgroup Ask that leader how many are in that subgroup Determine how many are in your subgroup. return \# in your subgroup + \# in other subgroup

## Recursive Counting - Divide and Conquer

RecursiveCount:
If you are the only person in group: return 1
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Recorsive Counting

- How long does this take with $n=2 \mathrm{~m}$ students?

Proof by induction:
You want to prove a statement $H(m)$ is true for all $m=0,1,2,3, \ldots$

StEps:

- Prove base case. Shaw $H(0)$ is true.
- Inductive Step.

Assume $H(m)$. Show $H(m+1)$ is true. " inductive hypothesis"

## Running Time - Proof by Induction

- Claim: For every number of students $n=2^{m}$

$$
\begin{aligned}
& T\left(2^{m}\right)=3 m+1 \\
& \text { Claim: } \begin{aligned}
& \text { Recursive Count takes } 3 m+1 \text { stops } \\
& \text { for an input } n=2^{m} \text {. That is } T\left(2^{m}\right)=3 m+1 \\
& \text { Know } T\left(2^{0}\right)=1 \quad \& T\left(2^{m}\right)=3+T\left(2^{m-1}\right) \\
& \text { Proof: } \\
& \text { Base cast } T\left(2^{\circ}\right)=1=3 \cdot 0+1 \quad \checkmark \\
& \text { Inductive case } \\
& T\left(2^{m+1}\right)=3+T\left(2^{m}\right) \\
&=3+3 m+1 \quad \text { (by inductive hypothesis) } \\
&=3(m+1)+1
\end{aligned}
\end{aligned}
$$

- In terms ot $\mathrm{n}, \quad T(n) \approx 3 \log _{2} n+1$

