

Definition: a **partial order** (sometimes called a *partially ordered set* or *poset*) is a pair, (D, \sqsubseteq) where:

D is a set and

\sqsubseteq is a reflexive, transitive, and anti-symmetric relation such that:

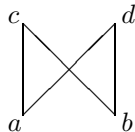
$$\forall x. x \sqsubseteq x$$

$$\forall x, y, z. (x \sqsubseteq y) \wedge (y \sqsubseteq z) \Rightarrow (x \sqsubseteq z)$$

$$\forall xy. (x \sqsubseteq y) \wedge (y \sqsubseteq x) \Rightarrow (x = y)$$

Definition: Let X be a subset of D , then:

$d \in D$ is an **upper-bound** (or UB) for X iff $\forall x \in X. (x \sqsubseteq d)$



c is an *upper-bound* for $\{a, b, c\}$ and
 c & d are both *upper-bounds* for $\{a, b\}$

Definition: d is a **least-upper-bound** (or LUB) for X iff:

(1) d is an *upper-bound* for X and

(2) if d' is any *upper-bound* for X then $d \sqsubseteq d'$

Property: if d & d' are both LUBs of X then $d = d'$

Proof: Since d and d' are LUBs of X , $(d \sqsubseteq d')$ and $(d' \sqsubseteq d)$, $(d = d')$ by *anti-symmetry* of \sqsubseteq

- If \sqsubseteq is not *anti-symmetric* we call it a *pre-order*

- We write $\sqcup X$ for the *least-upper-bound* of X , if it exists

What kinds of sets have LUBs ?

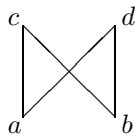
(1) no restriction: (*poset*)

(2) Every finite subset has a LUB (*lattice*)

(3) Every subset has a LUB (*complete lattice*)

Somewhere between 1 & 2 is our interesting class of sets

Definition: Given a *poset* (D, \sqsubseteq) , $X \subseteq D$ is **directed** iff every finite set $F \subseteq X$ has an *upper-bound* in X



directed	not directed
$\{a, c\}$	$\{a, b\}$
$\{a, b, c\}$	$\{a, b, c, d\}$
singleton sets	
pairs of \sqsubseteq	

Definition: a **complete-partial-order** (CPO) is a *poset* where every *directed* subset has a LUB .

Definition: a **pointed-CPO** is a CPO with a *least-element*, usually *bottom*, (\perp)

Any finite PO is a CPO since any directed set X is finite, choose $F = X$ so X has an *upper-bound* in X , that must be the LUB .

Interesting posets:

Singleton: $\{\cdot\}$, $\mathbf{1}$, or \mathbf{U}

2 element: $\begin{array}{c} \top \\ \vdots \\ \perp \end{array}$ or \mathbf{O}

3 element: $\begin{array}{ccc} & \text{true} & \text{false} \\ & \swarrow & \searrow \\ & \perp & \end{array}$, \mathbf{B} , or \mathbf{T}

$\omega^\top : (\mathbf{N} \cup \{\top\})$

If A is any set then $\langle \mathbf{P}(A), \subseteq \rangle$ is a CPO

$$X \subseteq \mathbf{P}(A), \bigsqcup X = \bigcup X$$

If A is any set then $\langle A, \{(x, y) \mid x = y\} \rangle$ is a CPO (but not *pointed*)

but, $\langle A \cup \{\perp\}, \{(x, y) \mid (x = y) \vee (x = \perp)\} \rangle$ is

Definition: Products of Posets

$$(D, \sqsubseteq_D) \times (E, \sqsubseteq_E) = (D \times E, \sqsubseteq_{D \times E})$$

where

$$(d, e) \sqsubseteq_{D \times E} (d', e') \text{ iff } (d \sqsubseteq_D) \wedge (e \sqsubseteq_E e')$$

If $P \subseteq D \times E$ is directed then P has a LUB

Claim: Let P be directed, $P_1 = \{x \mid \exists y.(x, y) \in P\}$ and $P_2 = \{y \mid \exists x.(x, y) \in P\}$, then P_1 & P_2 are directed

Proof: if $\{x_1, \dots, x_n\} \in P_1$ and $\{y_1, \dots, y_n\} \in P_2$ such that $\{(x_1, y_1), \dots, (x_n, y_n)\} \subseteq P$ then by the definition of $\sqsubseteq_{P_1 \times P_2}$, since P is directed so are P_1 and P_2

Partial Functions

If S and T are sets then $S \dashrightarrow T$ defines the set of **partial-functions** from S to T
 $f \sqsubseteq g$ iff $\forall x \in S$, if $f(x)$ is defined then so is $g(x)$, and $f(x) = g(x)$

Alternatively:

$$\text{graph}(f) = \{(x, y) \mid x \in S \wedge y = f(x)\}, \text{ then } f \sqsubseteq g \text{ iff } \text{graph}(f) \subseteq \text{graph}(g)$$

$X \subseteq [S \dashrightarrow T]$ is directed iff for any finite $F \subseteq X$, F has a UB in X

Let $X \subseteq S \dashrightarrow T$, $F \subseteq X$ with F finite.

Let $f \in F = \{(x_1, f(x_1)), \dots, (x_n, f(x_n))\}$

We want $g : S \rightarrow T$ such that $g(x) = f(x)$ if $\exists f \in X$ such that $f(x)$ is defined

Lemma I: if $f_1, f_2 \in X$ and $f_1(x)$ & $f_2(x)$ are defined, then $f_1(x) = f_2(x)$

Proof: $\{f_1, f_2\}$ is a finite subset of X , therefore it has an UB in X ,
 i.e. $\exists f_3$ such that $f_1 \sqsubseteq f_3$ and $f_2 \sqsubseteq f_3$. So $f_1(x) = f_3(x)$ and $f_2(x) = f_3(x)$

This all means that:

$S \dashrightarrow T$ is a CPO

Definition: Let $f : D \rightarrow E$ where $D \equiv (D, \sqsubseteq_D)$ and $E \equiv (E, \sqsubseteq_E)$
The function f is **monotone** iff $\forall d, d',$ if $d \sqsubseteq_D d'$ then $f(d) \sqsubseteq_E f(d')$

Lemma (Exchange):

Let D & D' be CPOs, $P \subseteq D$ & $Q \subseteq D'$ be directed, and D'' be any *poset*.
Let $f : D \times D \rightarrow D''$ be monotone, then:

$$\bigsqcup_{x \in P} \bigsqcup_{y \in Q} f(x, y) = \bigsqcup_{y \in Q} \bigsqcup_{x \in P} f(x, y)$$

Proof: Must show that:

$$\bigsqcup_{x \in P} \bigsqcup_{y \in Q} f(x, y) \sqsubseteq \bigsqcup_{y \in Q} \bigsqcup_{x \in P} f(x, y) \quad \text{and} \quad \bigsqcup_{y \in Q} \bigsqcup_{x \in P} f(x, y) \sqsubseteq \bigsqcup_{x \in P} \bigsqcup_{y \in Q} f(x, y)$$

First, $\bigsqcup_{x \in P}$ is a *LUB*, so it suffices to show that

$$\bigsqcup_{y \in Q} \bigsqcup_{x \in P} f(x, y) \text{ is an } \textit{upper-bound} \text{ for } \{ \bigsqcup_{y \in Q} f(x, y) \mid x \in P \}$$

So it remains to show that for any $x \in P$

$$\bigsqcup_{y \in Q} f(x, y) \sqsubseteq \bigsqcup_{y \in Q} \bigsqcup_{x \in P} f(x, y)$$

From there it remains to show that for any $y \in Q$

$$f(x, y) \sqsubseteq \bigsqcup_{y \in Q} \bigsqcup_{x \in P} f(x, y)$$

which is obvious since f is *monotone*

Options:

- (1) ... if these *LUBs* exist ...
- (2) ... let D'' be a CPO, then these *LUBs* exist ...

Lemma (Diagonal):

Let D & D' be CPOs, $P \subseteq D$ be directed, and $f : D \times D \rightarrow D'$ be monotone
then:

$$\bigsqcup_{x \in P} \bigsqcup_{y \in P} f(x, y) = \bigsqcup_{x \in P} f(x, x)$$

Proof: Homework Excercize #5

There are more notes here from the discussion the day after...

Definition: a function $f : D \rightarrow E$ where D & E are *pointed-CPOs* is **strict** iff it is *bottom-preserving*.

Definition: a function $f : D \rightarrow E$ where D & E are *CPOs* is **continuous** iff
 $\forall X \subseteq D$ where X is directed, $f(\bigsqcup X) = \bigsqcup_{x \in X} f(x)$.

This leads to the following properties:

(1) If f is *continuous*, then f is *monotone* (continuity \Rightarrow monotonicity)

Let $X = \{a, b\}$ where $a \sqsubseteq b$. Then X is *directed*.

Therefore $f(\bigsqcup X) = \bigsqcup \{f(x) \mid x \in X\}$, So

$f(b) = \bigsqcup \{f(a), f(b)\}$, therefore $f(a) \sqsubseteq f(b)$

(2) If f is *monotone* and D is *finite*, then f is *continuous*

(3) If f is *monotone* and D has no *infinite-increasing-chains*, then f is *continuous*

(4) f is *monotone* $\not\Rightarrow$ f is *continuous*

Example:

$f : \omega^\top \rightarrow \mathbf{O}$ with $f(\top) = \top$ and $f(n) = \perp$

$$\bigsqcup_{x \in X} f(x) = \perp \quad \text{and} \quad f\left(\bigsqcup_{x \in X} x\right) = \top$$

f is not *continuous* since $\bigsqcup f(x) \neq f(\bigsqcup x)$

Lemma: If $f : D \rightarrow E$ is continuous and $X \subseteq D$ is directed, then

$$\bigsqcup_{x \in X} f(x) \sqsubseteq f\left(\bigsqcup_{x \in X} x\right)$$

Proof: For any $x_0 \in X$, $x_0 \sqsubseteq \bigsqcup X$, since f is *monotone*, $f(x_0) \sqsubseteq \bigsqcup f(x)$

This is true for all $x_0 \in X$, so $\bigsqcup f(x)$ is an UB for $\{f(x) \mid x \in X\}$ and is *least*

Claim: If $f : D \rightarrow E$ and $g : E \rightarrow F$ are *continuous* then $(g \circ f) : D \rightarrow F$ is *continuous*

Proof: Must show that:

$$(g \circ f)\left(\bigsqcup_{x \in X} x\right) \sqsubseteq \bigsqcup_{x \in X} (g \circ f)(x)$$

Since X is directed, so is $\{f(x) \mid x \in X\}$ by the above Lemma.

$$\begin{aligned} g(f(\bigsqcup X)) &= g\left(\bigsqcup \{f(x) \mid x \in X\}\right) : \text{by continuity of } f \\ &= \bigsqcup \{g(f(x)) \mid x \in X\} : \text{by continuity of } g \\ &= \bigsqcup (g \circ f)(x) \end{aligned}$$

Definition: $[D \rightarrow E]$ is the set of all *continuous* functions from $D \rightarrow E$ such that

$$f \sqsubseteq_{D \rightarrow E} g \text{ iff } \forall d \in D. f(d) \sqsubseteq_E g(d)$$

For some S and T

$S \dashrightarrow T$ is a *partial-function-order*

$S \rightarrow T_\perp$ is a one to one correspondence by replacing *undefined* by \perp

Let $\varphi : (S \rightarrow T_\perp) \rightarrow (S \dashrightarrow T)$ be a function which converts a (total?) function from $S \rightarrow T_\perp$ to the corresponding partial function, such that $\varphi(f) = g$ where:

$$\begin{aligned} g(x) \text{ is undefined if } f(x) = \perp \\ g(x) = f(x) \text{ otherwise} \end{aligned}$$

Claim: $f \sqsubseteq_{S \rightarrow T_\perp} g$ iff $\varphi(f) \sqsubseteq \varphi(g)$

Proof: $\forall s \in S. (f(s) \sqsubseteq_{T_\perp} g(s)) \iff f(s) = g(s) \vee f(s) = \perp$
So, either $\varphi(f(s)) = g(s)$ or $\varphi(f(s))$ is undefined

Property: $[D \rightarrow E]$ is a CPO

Let $P \subseteq [D \rightarrow E]$ be *directed*

Define $g : D \rightarrow E$ such that $g(x) = \bigsqcup \{ f(x) \mid f \in P \}$

Claim: $g = \bigsqcup P$

Proof: We must show that:

- (1) g is *continuous*,
- (2) $\forall f \in P. f \sqsubseteq g$, and
- (3) $\forall f' \in \text{ub}(P). f' \sqsubseteq g$

?? I'm not sure where this section fits, my notes might be out of order...

Claim: If P is *directed* then $Q = \{ f(d) \mid f \in P \}$ is *directed*

Proof: Let $f_1, f_2 \in P$ with $f_1(d), f_2(d) \in Q$

If $f_1, f_2 \in P$ then that share an UB, say $f_3 \in P$ therefore $f_1(d)$ and $f_2(d)$ have an UB, $f_3(d) \in Q$.

So, Q is *directed* and g (from above?) is defined.

Let $Q \subseteq D$ be *directed*. We want to show that:

$$g\left(\bigsqcup_{x \in Q} x\right) = \bigsqcup_{x \in Q} g(x)$$

So:

$$\begin{aligned} g\left(\bigsqcup_{x \in Q} x\right) &= \bigsqcup_{f \in P} \left(f\left(\bigsqcup_{x \in Q} x\right)\right) && \text{by definition of } g \\ &= \bigsqcup_{f \in P} \left(\bigsqcup_{x \in Q} f(x)\right) && \text{since } f \text{ is continuous} \\ &= \bigsqcup_{x \in Q} \left(\bigsqcup_{f \in P} f(x)\right) && \text{by the Exchange lemma} \end{aligned}$$

Reminder:

- Continuous $\equiv f(\bigsqcup X) = \bigsqcup_{x \in X} f(x)$
- $\lambda x.x$ is *continuous*
 - $D \rightarrow E$, $\{(d, e_0) \mid d \in D\}$ is *continuous*
 - $E \rightarrow [D \rightarrow E]$ for any D, E is *continuous*

Definition: $K_{d,e} = \lambda e.\lambda d.e$

If $X \subseteq E$ is *directed* then:

$$K(\bigsqcup X) = \bigsqcup_{e \in X} K(e) = \bigsqcup_{e \in X} (\lambda d.e)$$

$$\bigsqcup_{e \in X} K(e)(d) = \bigsqcup_{e \in X} (e) = \bigsqcup X$$

$$\bigsqcup_{e \in X} \lambda e.\lambda d.e = \bigsqcup_{e \in X} (e) = \bigsqcup X$$

Definition: $S(f)(g)(x) = (f(x))(g(x))$

- $x : D$
- $g : D \rightarrow E$
- $f : D \rightarrow [E \rightarrow F]$
- $S : [[D \rightarrow [E \rightarrow F]] \rightarrow [D \rightarrow E] \rightarrow D] \rightarrow F$

Claim: S is *continuous*

Proof: Homework Excercize #6

Must show that if f & g are *continuous* then F is (*continuous* ?)

- Requires three results...

$$S : [[[D \Rightarrow [D \Rightarrow F]] \Rightarrow [D \Rightarrow E] \Rightarrow D] \Rightarrow F]$$

$$K : [D \Rightarrow [D \Rightarrow E]]$$

Application:

$$\frac{D \Rightarrow E \quad D}{E}$$

These are complete for propositional logic of *pure-implication*

Definition: Terms

- $t ::= x \mid K \mid S \mid (t_1 t_2)$
- $S = \lambda f.\lambda g.\lambda x.((f x)(g x))$
- $K = \lambda x.\lambda y.x$

Claim: Any lambda expression is equivalent to some combinatory term $[e]$

- (1) $[x] = x$
 - (2) $[e_1 e_2] = ([e_1] [e_2])$
 - (3) $[\lambda x.e] = [x]([e])$
- where $[x]$ is an operator on combinatory terms: (*Curry* bracket abstraction)
- $[x]x = ((S K) K) x$
 - $[x]t = (K t)$
 - $[x](t_1 t_2) = S([x]t_1)([x]t_2) = \lambda x.(([x]t_1)([x]t_2)) = \lambda x.(t_1 t_2)$

Conclusion... continuous functions are closed under composition